

Automatic Generation of Multi-Objective Polyhedral Compiler Transformations

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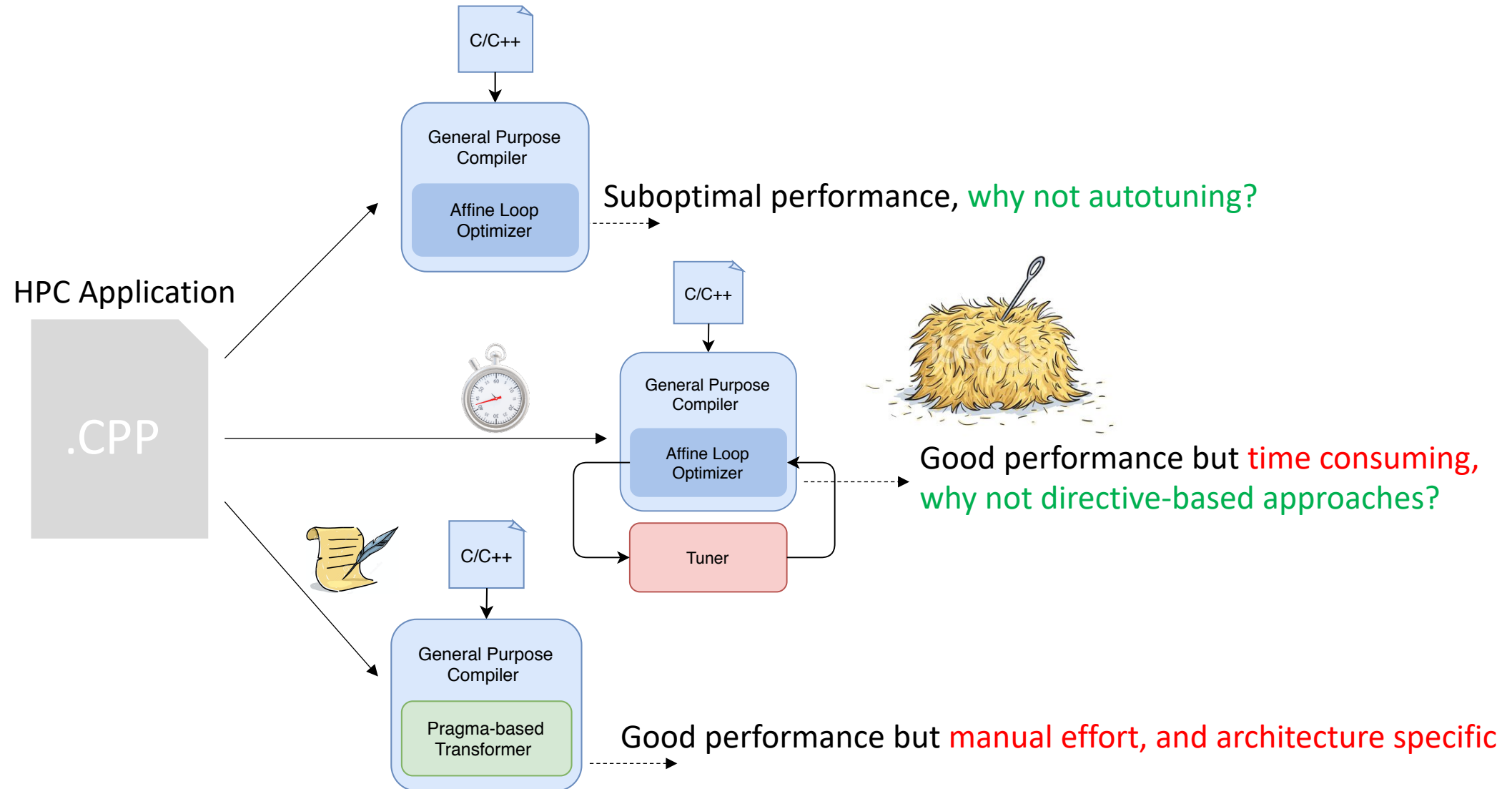
Talk Outline

- Challenges and Motivation
- Adaptive Scheduler
- Results
- Conclusion

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- **Challenges and Motivation**
- Adaptive Scheduler
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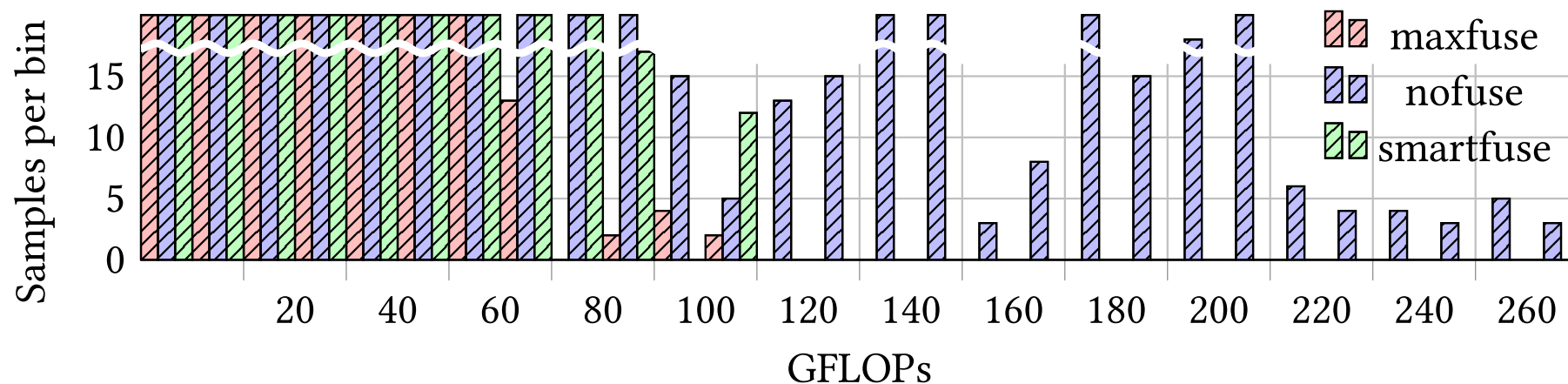
Challenges



Motivating Example

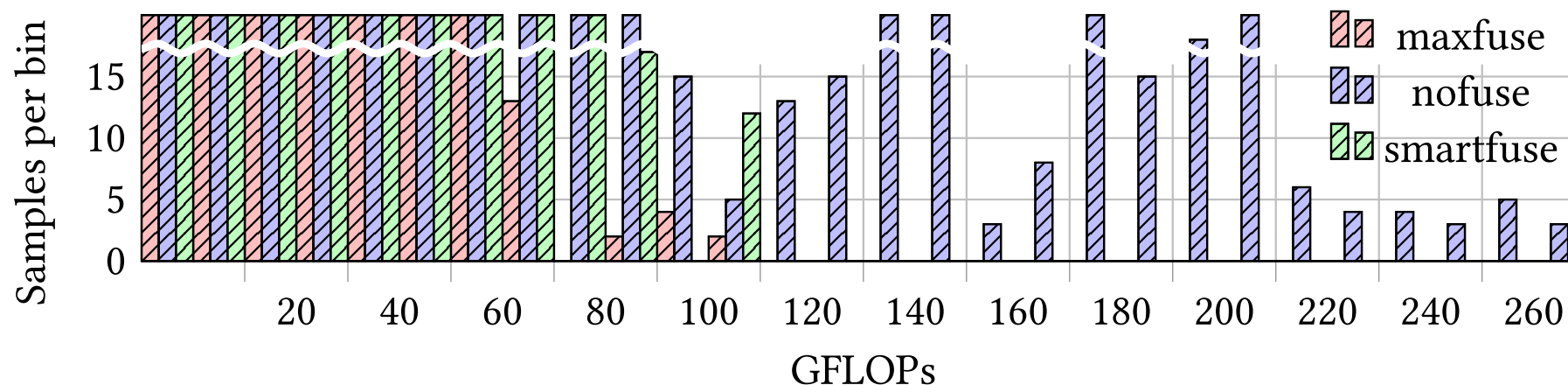
- Baseline: original and variants generated with the Pluto compiler
- More than 2000 variants
- Only 3/2000 variants reach 260 GFLOP/s
- MKL 0.75 TFLOP/s

```
for (int i = 0; i < N; ++i)
  for (int j = 0; j < N; ++j) {
S1:    tmp[i][j] = 0;
      for (int k = 0; k < N; ++k)
S2:    tmp[i][j] += alpha * A[i][k] * B[k][j];
      }
  for (int i = 0; i < N; ++i)
    for (int j = 0; j < N; ++j) {
S3:    D[i][j] *= beta;
      for (int k = 0; k < N; ++k)
S4:    D[i][j] += tmp[i][k] * C[k][j];
      }
  }
```

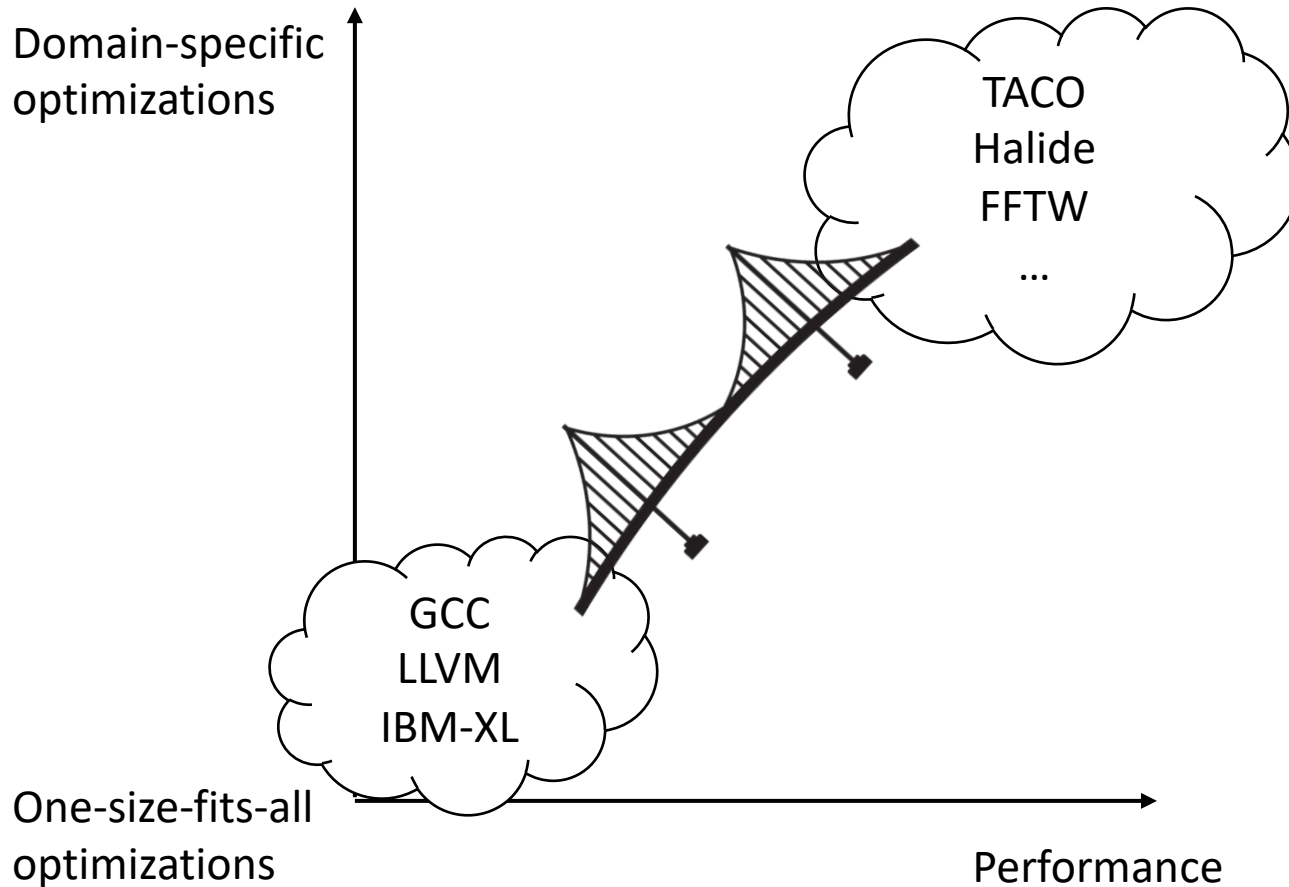


Motivating Example

- Baseline: original and variants generated with the Pluto compiler
 - More than 2000 variants
 - Only 3/2000 variants reach 260 GFLOP/s
 - MKL 0.75 TFLOP/s
- **Our work:**
 - Achieves almost the same performance (214 GFLOP/s)
 - By evaluating **only** five variants



Fundamental Contribution



This Work

- ✓ Bridge the gap between domain-specific and general-purpose optimizers
- ✓ Proposes a tractable approach to navigate large space of affine transformations
- ✓ Application and target customization

How to achieve this? Objective Functions

Objective Functions:

1. Outer Parallelism (OP)
2. Inner Parallelism (IP)
3. Outer Parallelism Inner Reuse (OPIR)
4. Stride Optimization (SO)

Stride Optimization:

```
for (int i = 0; i < NI; i++)  
  for (int j = 0; j < NJ; j++)  
    for (int k = 0; k < NK; k++) {  
  
      C[i][j] = 0;    // reuse  
  
      A[i][k] = 0;    // vectorization  
  
      B[k][j] = 0;    // penalty  
  
    }
```

Objective Functions an Example

```
for (int i = 0; i < N; ++i)
  for (int j = 0; j < N; ++j) {
S1:    tmp[i][j] = 0;
      for (int k = 0; k < N; ++k)
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S4:    D[i][j] += tmp[i][k] * C[k][j];
  }
```

```
// S1 and S2 are distributed as
// a by-product of the selected objectives.
```

```
// S1 /* ... */
      <OPIR, IP, SO>
for (int k = 0; k < N; ++k)    // reuse
  for (int i = 0; i < N; ++i)  // parallel
    for (int j = 0; j < N; ++j) // vectorizable
      tmp[i][j] += alpha * A[i][k] * B[k][j];
S2:
```

```
      <OP, IP, SO>
for (int i = 0; i < N; ++i)    // parallel
  for (int k = 0; k < N; ++k)  // reuse
    for (int j = 0; j < N; ++j) // vectorizable
      tmp[i][j] += alpha * A[i][k] * B[k][j];
S2:
```

```
      <IP, OPIR, OP>
for (int j = 0; j < N; ++j)    // vectorizable
  for (int k = 0; k < N; ++k)  // reuse
    for (int i = 0; i < N; ++i) // parallel
      tmp[i][j] += alpha * A[i][k] * B[k][j];
S2:
```

Objective Functions an Example

```
// S1 and S2 are distributed as  
// a by-product of the selected objectives.
```

```
// S1 /* ... */  
        <OPIR, IP, SO>
```

```
for (int k = 0; k < N; ++k) // reuse
```

Order of Objectives matter!

Large exploration space!

```
S2: tmp[i][j] += alpha * A[i][k] * B[k][j],  
}  
for (int i = 0; i < N; ++i)
```

```
for (int k = 0; k < N; ++k) // reuse  
    for (int j = 0; j < N; ++j) // vectorizable  
        tmp[i][j] += alpha * A[i][k] * B[k][j].  
S2:
```

How to decide the best sequence of Objective Functions?

```
S4: D[i][j] += tmp[i][k] * C[k][j].
```

```
for (int j = 0; j < N; ++j) // vectorizable
```

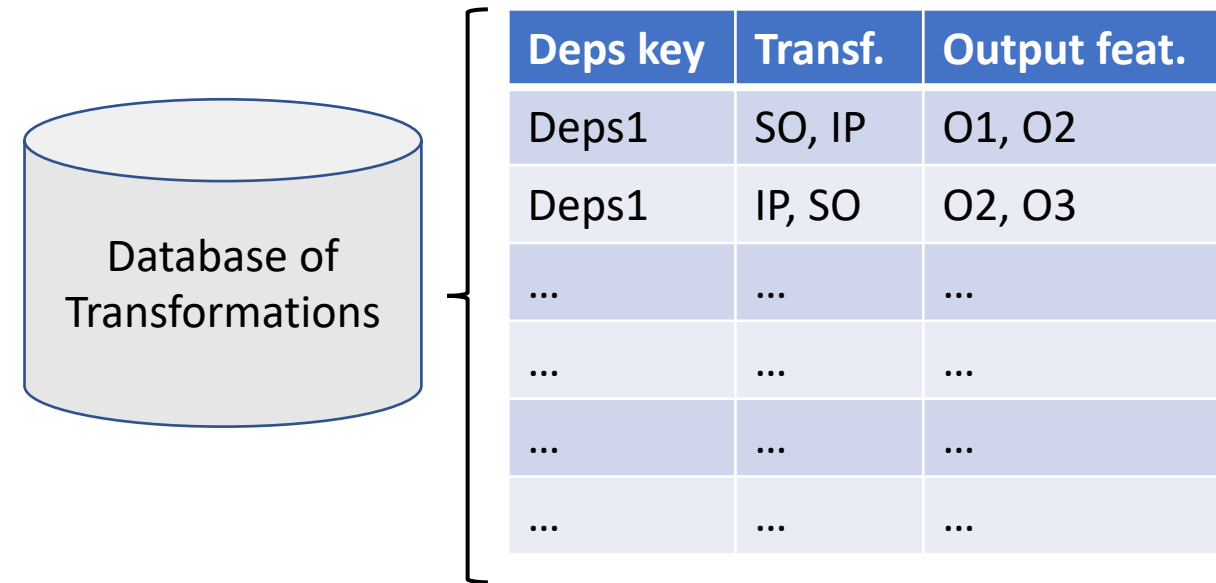
How to characterize, traverse and quantify such large loop transformation spaces?

Talk Outline

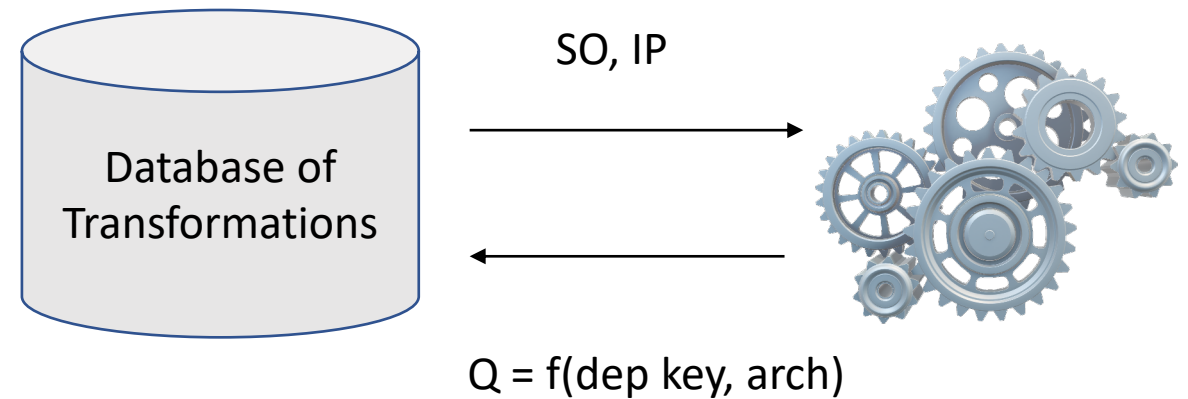
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Adaptive Scheduler: Two components

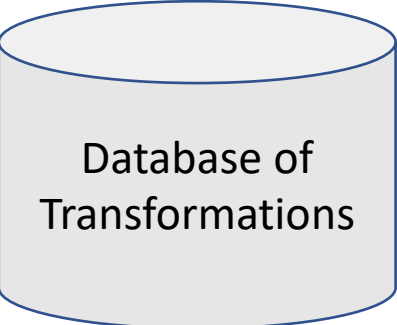
- Offline Stage: Building a Database of Transformations



- Online Stage: Select transformations



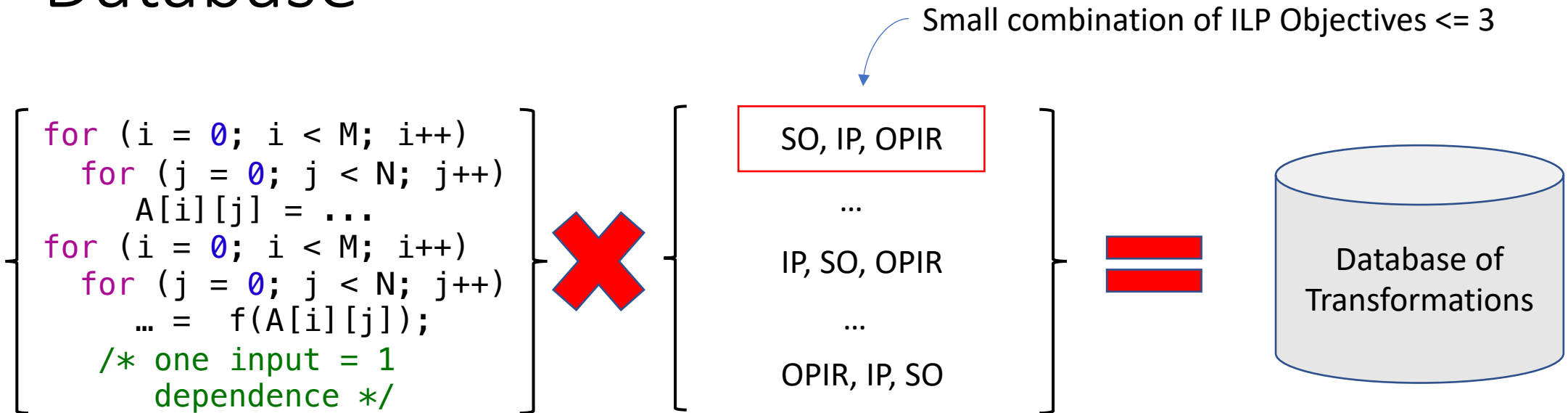
Database of Transformations



Database of Transformations

Dependences key	Transformations	Output features
Dependence 1	Stride Opt., Inner Paral.	O1, O2
Dependence 1	Inner Paral., Stride Opt.	O2, O3
• Model a single dependence polyhedron • Extremely simple to allow scalability • Dependences keys encompass structural information and dependence features (e.g. 1-to-1 or 1-to-N)	• Each transformation implemented as an Integer Linear Program (ILP) cost function • Characterize combined behavior of transformations • Order matters	• Capture transformation results and properties (e.g. number of parallel loops, total number of loops, stride access, and more) • Used to compute a score for each candidate transformation

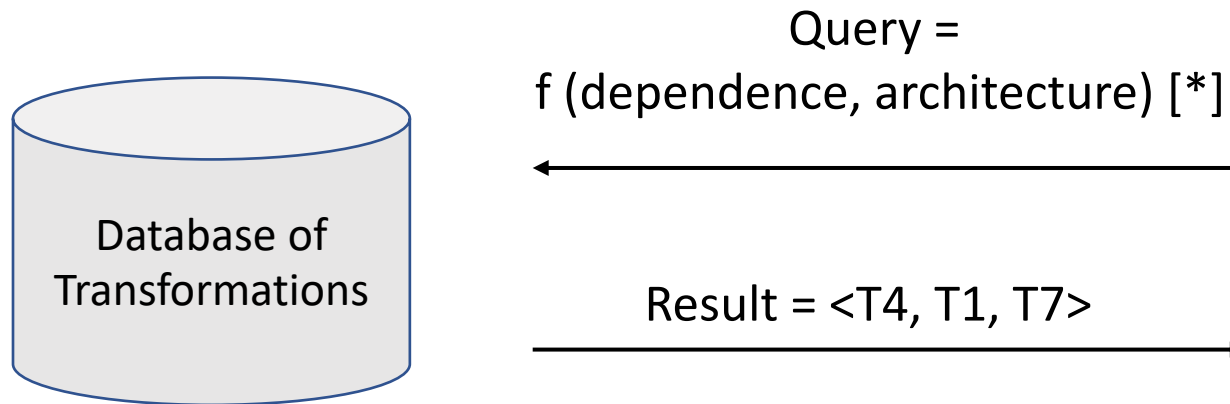
Offline Stage: Building the Transformation Database



Synthesized or extracted from
larger representative
applications

Candidate transformation is a tuple of ILP
objectives, e.g. all 3 permutations in a set of
10 transformations = 720 transformation
candidates

Online Stage: Selecting Transformations



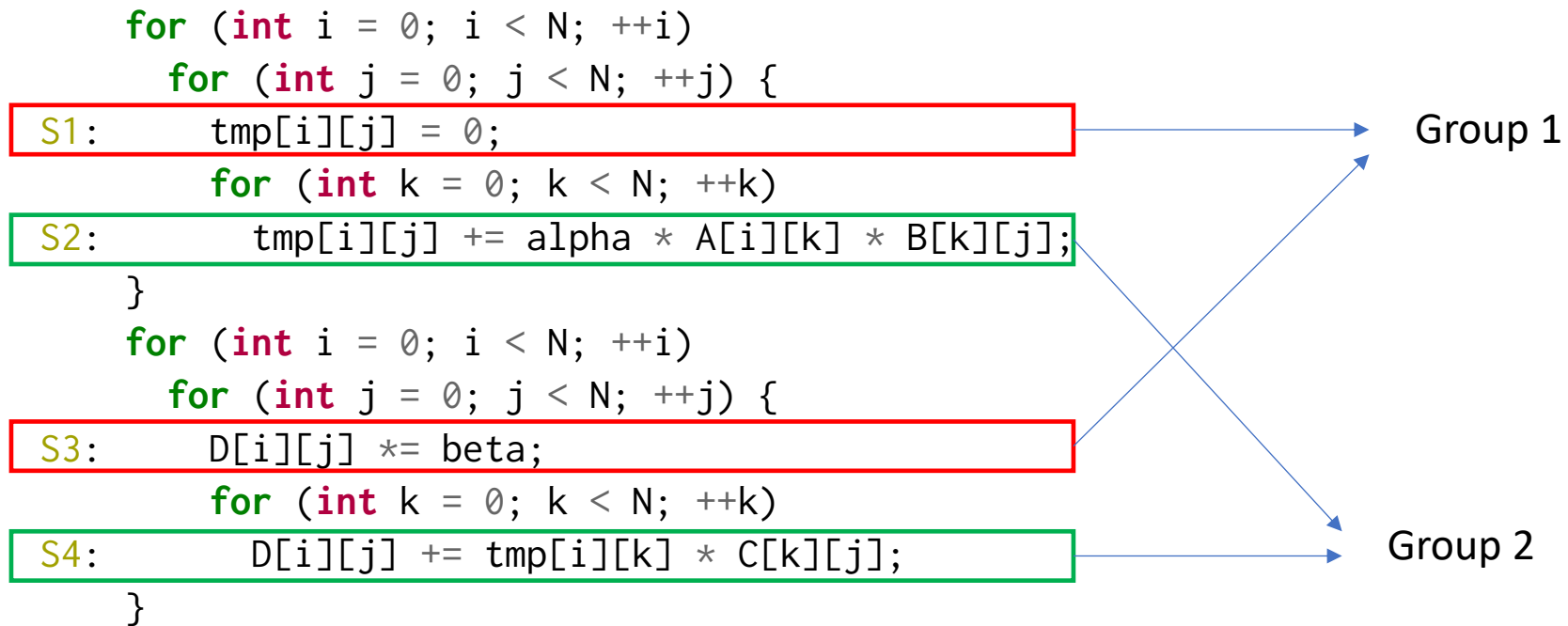
$$\text{Transformation} = \text{Query 1} + \text{Query 2}$$

Diagram illustrating the transformation selection process. The transformation is composed of two parts: Query 1 (represented by the set $\langle T2, T5, T4 \rangle$) and Query 2 (represented by the set $\langle T4, T1, T7 \rangle$). Red circles and arrows highlight the mapping from the queries to their respective transformation sets.

1. Divide statements into equivalence classes with 5 partition policies
2. Select most influential dependences in each partition

A step-by-step Example

- Step 1: Partition statements



A step-by-step Example

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      }
}
```

- Step 1: Partition statements
- Step 2: Ranking dependences

D1: S1 $\rightarrow$ S2
D2: S2 $\rightarrow$ S2
D3: S2 $\rightarrow$ S4
D4: S3 $\rightarrow$ S4
D5: S4 $\rightarrow$ S4

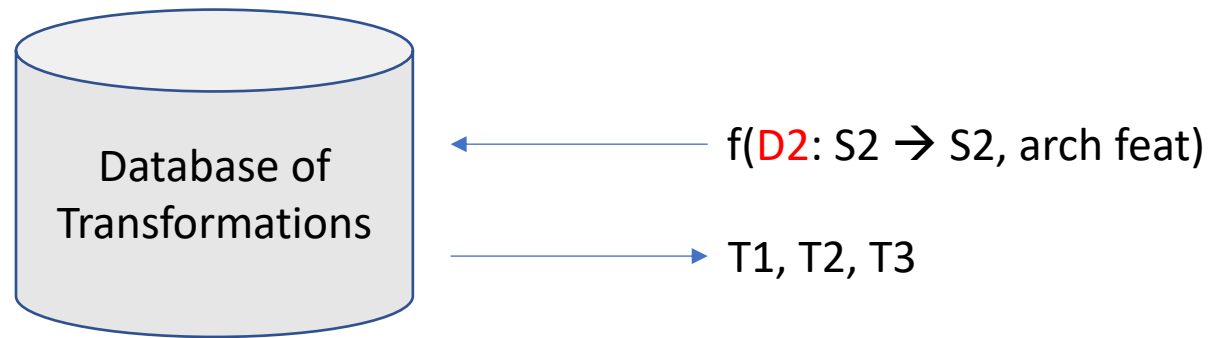


D2: S2 $\rightarrow$ S2
D5: S4 $\rightarrow$ S4
D3: S2 $\rightarrow$ S4
D4: S3 $\rightarrow$ S4
D1: S1 $\rightarrow$ S2

A step-by-step Example

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S4:    D[i][j] += tmp[i][k] * C[k][j];
    }
```

- Step 1: Partition statements
- Step 2: Ranking dependences
- Step 3: Query Database



D5 already covered by D2

A step-by-step Example

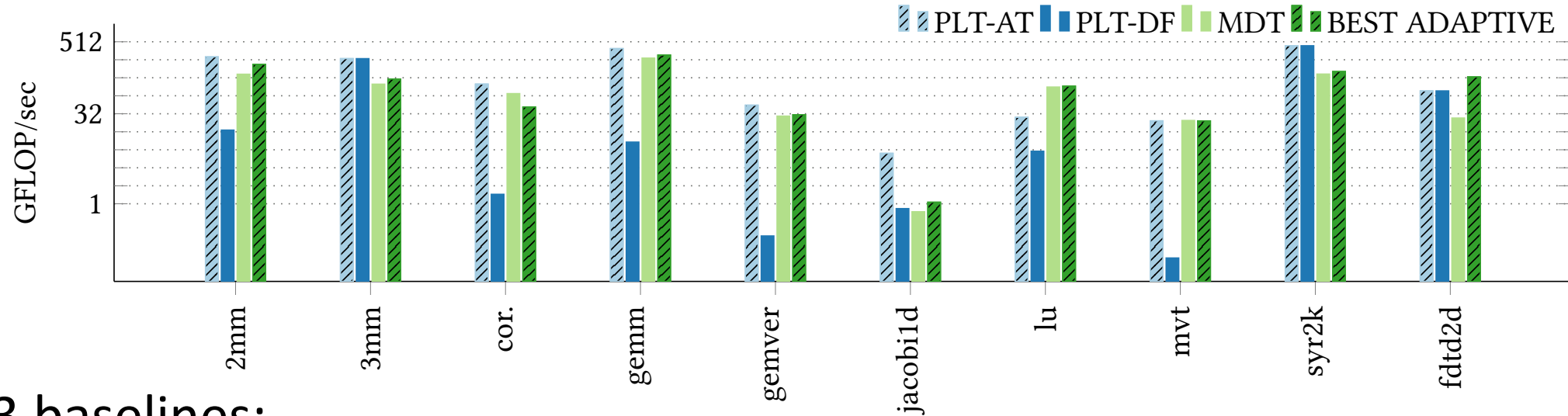
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```

- Step 1: Partition statements
- Step 2: Ranking dependences
- Step 3: Query Database
- Step 4: Compute lexicographic optimum and apply transformations

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Results... Some of them



- 3 baselines:

- Pluto autotune (PLT-AT): Best of thousand program variants
- Pluto Default (PLT-DF): Pluto out-of-the-box
- Model driven transformations (MDT): Single optimized variant using predefined order of transformations

- Best adaptive: Best variants for the 5 proposed transformation policies
- Machine: 48-core Intel Xeon Platinum (MKL 0.75 TFLOP/s)

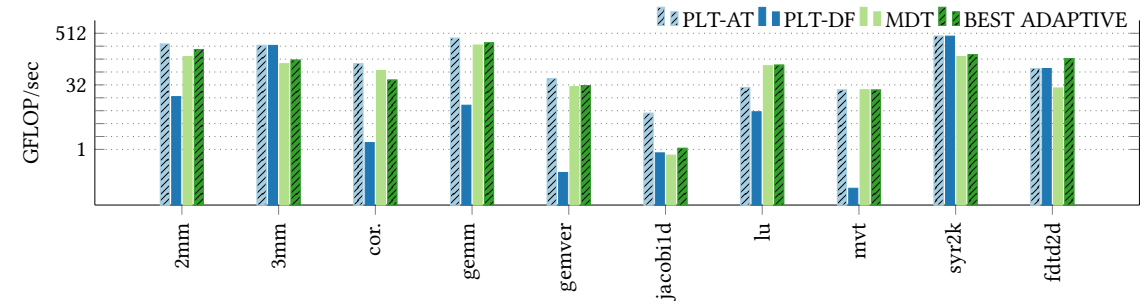
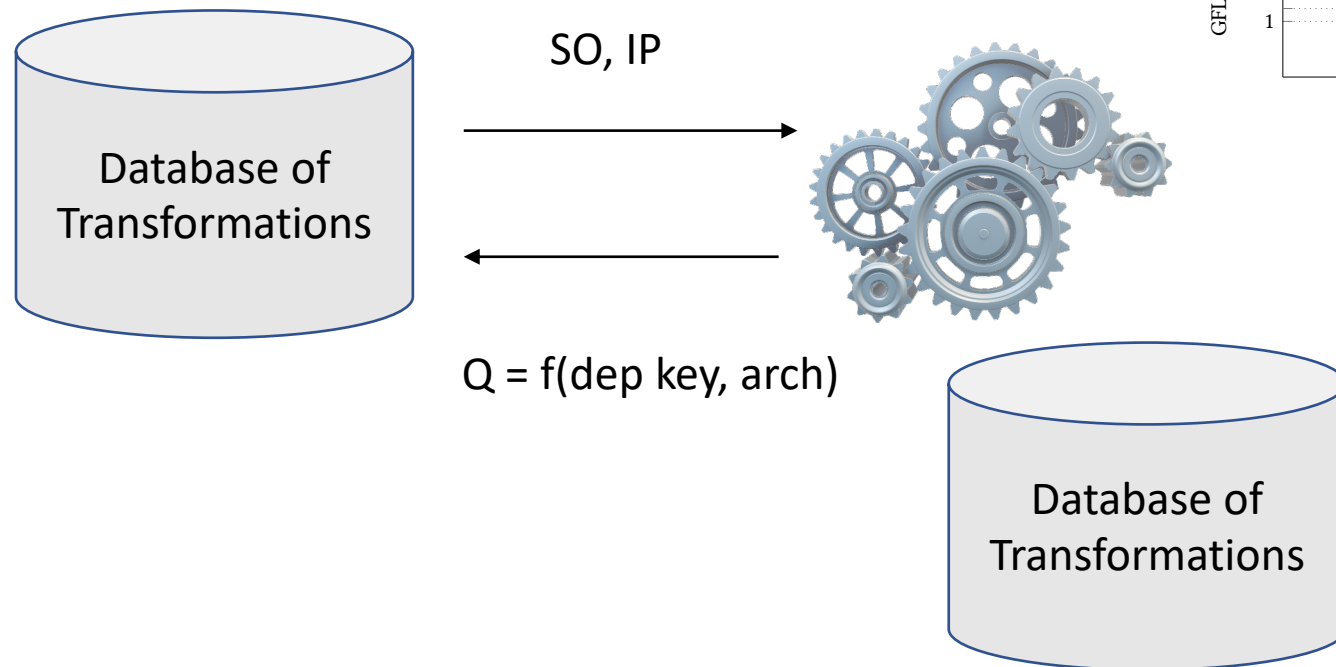
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Conclusion

- Scalable approach to traverse large transformation space
- Order of objectives is not hardcoded
- Reduction of tuning time
- Competitive performance with minimal exploration

Our adaptive scheduler bridges the gap between general-purpose and domain-specific compilation flow



Deps key	Transf.	Output feat.
Deps1	SO, IP	O1, O2
Deps1	IP, SO	O2, O3
...	...	...